## A note on the spatiality of localic products of countably based sober spaces

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Let **Top** be the category of topological spaces and continuous maps, **Loc** the category of locales, and  $\Omega: \mathbf{Top} \to \mathbf{Loc}$  the usual functor mapping spaces to locales.  $\Omega$  preserves colimits (since it has a right adjoint  $\mathbf{pt}: \mathbf{Loc} \to \mathbf{Top}$ ), but  $\Omega$  does not preserve finite products in general. The purpose of this note is to investigate subcategories of countably based sober spaces for which the restriction of  $\Omega$  does preserve finite products.

Let  $S_0$  be the countable space defined in [2]. The underlying set of  $S_0$  is  $\mathbb{N}^{<\mathbb{N}}$ , the set of finite sequences of natural numbers. A subbasis for the open subsets of  $S_0$  is given by sets of the form  $\{\tau \in \mathbb{N}^{<\mathbb{N}} \mid \sigma \not\preceq \tau\}$ , where  $\sigma \in \mathbb{N}^{<\mathbb{N}}$  and  $\preceq$  is the prefix relation. Note that  $S_0$  is a countably based sober space, its specialization order is the reverse of the prefix relation, and that  $S_0$  has uncountably many distinct open sets ([2]; Proposition 6.1).

We first show that the localic product  $\Omega(S_0) \times_l \Omega(S_0)$  is not spatial by describing a winning strategy for Player I in the game  $\mathcal{G}(S_0, S_0)$  defined by T. Plewe (see Theorem 1.1 in [5] and the paragraph above it for a definition of the game). The proof strategy for the following lemma is essentially the same as P. Johnstone's proof that  $\Omega(\mathbb{Q}) \times_l \Omega(\mathbb{Q})$  is not spatial (see Proposition II-2.14 of [4]), but the game theoretic approach allows us to hide the use of transfinite ordinals.

**Lemma 1.** The localic product  $\Omega(S_0) \times_l \Omega(S_0)$  is not spatial.

*Proof.* We denote the length of  $\sigma \in \mathbb{N}^{<\mathbb{N}}$  by  $|\sigma|$ . The empty string is denoted as  $\varepsilon$ , and the string consisting of m zeros is written  $0^{(m)}$ . The string obtained by appending  $n \in \mathbb{N}$  to  $\sigma \in \mathbb{N}^{<\mathbb{N}}$  is written  $\sigma \diamond n$ . We also write  $\sigma \diamond \tau$  for the concatenation of strings. For  $\sigma, \tau \in \mathbb{N}^{<\mathbb{N}}$ , define

$$F_{\sigma,\tau} = \{s \diamond n \mid n \in \mathbb{N} \& s \in \mathbb{N}^{<\mathbb{N}} \& s \preceq \sigma \& s \diamond n \not\preceq \sigma \& n \leq |\sigma| + |\tau|\}, \text{ and} \\ U_{\sigma,\tau} = \{t \in \mathbb{N}^{<\mathbb{N}} \mid (\forall s \in F_{\sigma,\tau}) s \not\preceq t\}.$$

Then  $\sigma \in U_{\sigma,\tau}$ , hence  $\mathcal{U} = \{ U_{\sigma,\tau} \times U_{\tau,\sigma} \mid \sigma, \tau \in \mathbb{N}^{<\mathbb{N}} \}$  is an open cover of  $S_0 \times S_0$ .

Observe that if  $s \in U_{\sigma,\tau}$  and every element of s is less than or equal to  $|\sigma| + |\tau|$  then  $s \leq \sigma$ . Also note that if  $U \subseteq S_0$  is open and  $\sigma \in U$ , then there exist infinitely many  $n \in \mathbb{N}$  such that every string that has  $\sigma \diamond n$  as a prefix is also in U. This is because there must be a finite  $F \subseteq S_0$ such that the basic open set  $W = \{\tau \in \mathbb{N}^{\leq \mathbb{N}} \mid (\forall s \in F) \ s \not \leq \tau\}$  satisfies  $\sigma \in W \subseteq U$ . Fix any  $n \in \mathbb{N}$  that is strictly larger than any element contained in any of the strings in F. Then for each  $s \in F$ , we have that  $s \not \leq \sigma \diamond n$  and  $\sigma \diamond n \not \leq s$ , hence no extension of  $\sigma \diamond n$  has s as a prefix. It follows that every extension of  $\sigma \diamond n$  must be in W and therefore also in U.

Player I initializes the game by playing  $S_0 \times S_0$  and the open covering  $\mathcal{U}$  of  $S_0 \times S_0$ . The game begins with round 1. For convenience, define  $V_0 = W_0 = S_0$ , and  $x_0 = y_0 = \varepsilon$ , and  $m_0 = n_0 = 0$ . Player I's strategy for the *i*-th round  $(i \ge 1)$  proceeds as follows:

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- 1. Player I chooses  $m_i \in \mathbb{N}$  such that every sequence extending  $y_{i-1} \diamond m_i$  is in  $W_{i-1}$ . Player I then plays  $x_i = x_{i-1} \diamond n_{i-1} \diamond 0^{(m_i)}$ .
- 2. Player II must respond with an open subset  $V_i \subseteq V_{i-1}$  containing  $x_i$ .
- 3. Next, Player I finds distinct  $n_i$  and  $n'_i$  in  $\mathbb{N}$  such that any sequence that has either  $x_i \diamond n_i$  or  $x_i \diamond n'_i$  as a prefix is in  $V_i$ . Player I plays  $y_i = y_{i-1} \diamond m_i \diamond 0^{(n_i+n'_i)}$ .
- 4. Player II must respond with an open subset  $W_i \subseteq W_{i-1}$  containing  $y_i$ .

The game then continues on to round i + 1.

We show that at the end of each round  $i \geq 1$ , the open rectangle  $V_i \times W_i$  chosen by Player II is not a subset of any open rectangle in  $\mathcal{U}$ . Fix any  $\sigma, \tau \in \mathbb{N}^{<\mathbb{N}}$  with  $x_i \in U_{\sigma,\tau}$  and  $y_i \in U_{\tau,\sigma}$ . Since  $y_{i-1} \leq y_i$  we have  $y_{i-1} \in U_{\tau,\sigma}$ , and an inductive argument (keep reading) yields  $y_{i-1} \leq \tau$ . Using the fact that  $|y_{i-1}| \geq n_{i-1}$  it can be shown that every element occurring in  $x_i$  is less than or equal to  $|y_{i-1}| \leq |\sigma| + |\tau|$ , hence the assumption  $x_i \in U_{\sigma,\tau}$  and the observation at the top of the second paragraph of this proof implies  $x_i \leq \sigma$ . Similarly, every element of  $y_i$  is less than or equal to  $|x_i| \leq |\tau| + |\sigma|$ , hence  $y_i \in U_{\tau,\sigma}$  implies  $y_i \leq \tau$  (thereby completing the inductive argument). Either  $x_i \diamond n_i \not\preceq \sigma$  or  $x_i \diamond n'_i \not\preceq \sigma$ , and  $n_i, n'_i \leq |y_i| \leq |\sigma| + |\tau|$ , thus  $x_i \diamond n_i \not\in U_{\sigma,\tau}$  or  $x_i \diamond n'_i \not\in U_{\sigma,\tau}$ , but both  $x_i \diamond n_i$  and  $x_i \diamond n'_i$  are in  $V_i$ , so we conclude that  $V_i \times W_i \not\subseteq U_{\sigma,\tau} \times U_{\tau,\sigma}$ . Therefore, the above strategy is winning for Player I, hence  $\Omega(S_0) \times_l \Omega(S_0)$  is not spatial.  $\Box$ 

Let  $\omega$ **Sob** be the category of countably based sober spaces, and let **QPol** be the category of quasi-Polish spaces [1].  $\omega$ **Sob** and **QPol** are closed under countable limits (as defined in **Top**), and the restriction of  $\Omega$  to **QPol** preserves all countable limits (Theorems 4.4 and 4.5 of [3]).

**Theorem 1.** Assume C is a full subcategory of  $\omega$ **Sob** satisfying:

(1) C is closed under finite limits (as defined in **Top**),

- (2) the restriction of  $\Omega$  to  $\mathcal{C}$  preserves finite products,
- (3)  $\mathcal{C}$  contains  $\mathcal{P}(\mathbb{N})$  (the powerset of  $\mathbb{N}$  with the Scott-topology), and
- (4) every space in  $\mathcal{C}$  is co-analytic (i.e., homeomorphic to a  $\Pi_1^1$ -subspace of  $\mathcal{P}(\mathbb{N})$ ).

Then C is a full subcategory of **QPol**.

Proof. Assume for a contradiction that there is some space X in  $\mathcal{C}$  which is not quasi-Polish. X is co-analytic by (4), hence Theorem 7.2 of [2] implies there is a  $\Pi_2^0$ -subspace Y of X which is homeomorphic to either  $S_0$  or  $\mathbb{Q}$  (the candidates  $S_D$  and  $S_1$  mentioned in [2] can be omitted because X is sober). Since Y is a  $\Pi_2^0$ -subspace of X, it is the equalizer of a pair of continuous functions  $f, g: X \to \mathcal{P}(\mathbb{N})$  (see the concluding section of [1]), hence Y is in  $\mathcal{C}$  by (1), (3), and the assumption that  $\mathcal{C}$  is a full subcategory of  $\omega$ **Sob**. Lemma 1 and the fact that  $\Omega(\mathbb{Q}) \times_l \Omega(\mathbb{Q})$ is not spatial imply  $\Omega(Y \times Y) \neq \Omega(Y) \times_l \Omega(Y)$ . Therefore, the restriction of  $\Omega$  to  $\mathcal{C}$  does not preserve products, which contradicts (2).

We conjecture that the above theorem still holds if (3) is omitted. It is consistent with ZFC to replace "co-analytic" in (4) with any level of the projective hierarchy, and we conjecture that it is consistent with ZF+(Dependent Choice) if (4) is removed completely.

## References

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