

Continuity in constructive analysis

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Aim:

Constructive analysis, with constructions \sim good algorithms.

Errett Bishop 1967: "Foundations of Constructive Analysis"

The modulus of continuity ω is an indispensable part of the definition of a continuous function on a compact interval, although sometimes it is not mentioned explicitly. In the same way, the moduli of continuity of the restrictions of f to each compact subinterval are indispensable parts of the definition of a continuous function f on a general interval.

A **continuous function** $f: (X, \rho, Q) \rightarrow (Y, \sigma, R)$ for separable metric spaces is given by

$$h: Q \rightarrow \mathbb{N} \rightarrow R \quad \text{approximating map}$$

plus $\alpha, \omega, \gamma, \delta$ depending on w, r (center and radius of a ball):

- ▶ $\alpha: Q \rightarrow \mathbb{Z}^+ \rightarrow \mathbb{Z}^+ \rightarrow \mathbb{N}$ such that $(h(u, n))_n$ (for $\rho(u, w) \leq \frac{1}{2r}$) is a **Cauchy sequence** with **modulus** $\alpha_{w,r}(p)$;
- ▶ a **modulus** $\omega: Q \rightarrow \mathbb{Z}^+ \rightarrow \mathbb{Z}^+ \rightarrow \mathbb{Z}^+$ of (uniform) continuity, such that for $n \geq \alpha_{w,r}(p)$ and $\rho(u, w), \rho(v, w) \leq \frac{1}{2r}$

$$\rho(u, v) \leq \frac{2}{2^{\omega_{w,r}(p)}} \rightarrow \sigma(h(u, n), h(v, n)) \leq \frac{1}{2^p};$$

- ▶ maps $\gamma: Q \rightarrow \mathbb{Z}^+ \rightarrow R, \delta: Q \rightarrow \mathbb{Z}^+ \rightarrow \mathbb{Z}^+$ such that $\gamma(w, r)$ and $\delta(w, r)$ are center and radius of a ball containing all $h(u, n)$ (for $\rho(u, w) \leq \frac{1}{2r}$):

$$\rho(u, w) \leq \frac{1}{2r} \rightarrow \sigma(h(u, n), \gamma(w, r)) \leq \frac{1}{2^{\delta(w, r)}}.$$

$\alpha, \omega, \gamma, \delta$ are required to have monotonicity properties.

f given by **type-1 data** only.

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Example: Inverse map $(0, \infty) \rightarrow \mathbb{R}$

Let $0 < c < d$, and q be minimal such that $\frac{1}{2^q} \leq c$. Then inv is given by

- ▶ the approximating map $h(a, n) := \frac{1}{a}$
- ▶ the Cauchy modulus $\alpha(c, d, p) := 0$
- ▶ the modulus $\omega(c, d, p) := p + 2q + 1$ of uniform continuity, for

$$|a - b| \leq \frac{1}{2^{p+2q}} \rightarrow \left| \frac{1}{a} - \frac{1}{b} \right| = \left| \frac{b - a}{ab} \right| \leq \frac{1}{2^p},$$

because $ab \geq \frac{1}{2^{2q}}$

- ▶ the center $\gamma(c, d) := \frac{c}{c^2 - d^2}$ and radius $\delta(c, d) := \frac{d}{c^2 - d^2}$ of a ball containing all $\frac{1}{a}$ for $|a - c| \leq d$.

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- ▶ **Application** $f(x)$ must (and can) be defined separately, since the approximating map operates on approximations only.
- ▶ $f(x)$ is independent from w, r .
- ▶ Application is compatible with equality on real numbers:

$$x = y \rightarrow f(x) = f(y).$$

- ▶ f has ω as a modulus of uniform continuity:

$$|x - y| \leq \frac{1}{2^{\omega(\rho)}} \rightarrow |f(x) - f(y)| \leq \frac{1}{2^{\rho}}.$$

- ▶ **Composition** can be defined.

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Algorithms in constructive proofs?

Theorem. Every totally bounded set $A \subseteq \mathbb{R}$ has an infimum y .

Proof.

Given $\varepsilon = \frac{1}{2^p}$, let $a_0 < a_1 < \dots < a_{n-1}$ be an ε -net:

$\forall x \in A \exists i < n (|x - a_i| < \varepsilon)$. Let $b_p = \min\{a_i \mid i < n\}$. $y := \lim_p b_p$. \square

Corollary. $\inf_{x \in [a,b]} f(x)$ exists, for $f: [a, b] \rightarrow \mathbb{R}$ continuous.

Proof.

Given ε , pick $a = a_0 < a_1 < \dots < a_{n-1} = b$ s.t. $a_{i+1} - a_i < \omega(\varepsilon)$.

Then $f(a_0), f(a_1), \dots, f(a_{n-1})$ is an ε -net for f 's range. \square

Many $f(a_i)$ need to be computed.

Aim: Get x with $f(x) = \inf_{y \in [a,b]} f(y)$ and a better algorithm, assuming convexity.

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Given $\varepsilon = \frac{1}{2^p}$, let $a_0 < a_1 < \dots < a_{n-1}$ be an ε -net:

$\forall x \in A \exists i < n (|x - a_i| < \varepsilon)$. Let $b_p = \min\{a_i \mid i < n\}$. $y := \lim_p b_p$. \square

Corollary. $\inf_{x \in [a,b]} f(x)$ exists, for $f: [a, b] \rightarrow \mathbb{R}$ continuous.

Proof.

Given ε , pick $a = a_0 < a_1 < \dots < a_{n-1} = b$ s.t. $a_{i+1} - a_i < \omega(\varepsilon)$.

Then $f(a_0), f(a_1), \dots, f(a_{n-1})$ is an ε -net for f 's range. \square

Many $f(a_i)$ need to be computed.

Aim: Get x with $f(x) = \inf_{y \in [a,b]} f(y)$ and a better algorithm, assuming convexity.

Algorithms in constructive proofs?

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Intermediate value theorem

Let $a < b$ be rationals. If $f: [a, b] \rightarrow \mathbb{R}$ is continuous with $f(a) \leq 0 \leq f(b)$, and with a uniform modulus of increase

$$\frac{1}{2^p} < d - c \rightarrow \frac{1}{2^{p+q}} < f(d) - f(c),$$

then we can find $x \in [a, b]$ such that $f(x) = 0$.

Proof (trisection method).

1. **Approximate Splitting Principle.** Let x, y, z be given with $x < y$. Then $z \leq y$ or $x \leq z$.
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Extracted term

```
[k0]
left((cDC rat@@rat)(1@2)
  ([n1]
    (cId rat@@rat=>rat@@rat)
    ([cd3]
      [let cd4
        ((2#3)*left cd3+(1#3)*right cd3@
          (1#3)*left cd3+(2#3)*right cd3)
        [if (0<=(left cd4*left cd4-2+
          (right cd4*right cd4-2)))/2)
          (left cd3@right cd4)
          (left cd4@right cd3)]]))
(IntToNat(2*k0)))
```

where `cDC` is a form of the recursion operator.

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Express this view as **invariance under realizability** axioms

$$\text{Inv}_A: A \leftrightarrow \exists_z (z \mathbf{r} A).$$

Consequences are **choice** and **independence of premise** (Troelstra):

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Derivatives

Let $f, g: I \rightarrow \mathbb{R}$ be continuous. g is called **derivative** of f with modulus $\delta_f: \mathbb{Z}^+ \rightarrow \mathbb{N}$ of differentiability if for $x, y \in I$ with $x < y$,

$$y \leq x + \frac{1}{2^{\delta_f(p)}} \rightarrow |f(y) - f(x) - g(x)(y - x)| \leq \frac{1}{2^p}(y - x).$$

A bound on the derivative of f serves as a Lipschitz constant of f :

Lemma (BoundSlope)

Let $f: I \rightarrow \mathbb{R}$ be continuous with derivative f' . Assume that f' is bounded by M on I . Then for $x, y \in I$ with $x < y$,

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Infimum of a convex function

Let $f, f' : [a, b] \rightarrow \mathbb{R}$ ($a < b$) be continuous and f' derivative of f . Assume that f is **strictly convex** with witness q , in the sense that $f'(a) < 0 < f'(b)$ and

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Then we can find $x \in (a, b)$ such that $f(x) = \inf_{y \in [a, b]} f(y)$.

Proof.

- ▶ To obtain x , apply the intermediate value theorem to f' .
- ▶ To prove $\forall_{y \in [a, b]} (f(x) \leq f(y))$ (this is “non-computational”, i.e., a Harrop formula) one can use the standard arguments in classical analysis (Rolle’s theorem, mean value theorem). \square

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Conclusion

Aim: constructive analysis, with constructions \sim good algorithms.
Then extract these algorithms from proofs (realizability).

- ▶ Use order locatedness: given $c < d$, for all u

$$\forall_{v \in V}(c \leq \rho(u, v)) \vee \exists_{v \in V}(\rho(u, v) \leq d).$$

- ▶ Avoid total boundedness (existence of ε -nets).

Generally

- ▶ View constructive analysis as an **extension** of classical analysis.
- ▶ Formalize proofs in **TCF** (based on the Scott-Ershov model of partial continuous functionals), extract algorithms (in Minlog).
- ▶ **Data** are important (real number, continuous function ...).
- ▶ **Low type levels**: continuous $f: \mathbb{R} \rightarrow \mathbb{R}$ determined by its values on the rationals \mathbb{Q} .

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Then extract these algorithms from proofs (realizability).

- ▶ Use order locatedness: given $c < d$, for all u

$$\forall_{v \in V}(c \leq \rho(u, v)) \vee \exists_{v \in V}(\rho(u, v) \leq d).$$

- ▶ Avoid total boundedness (existence of ε -nets).

Generally

- ▶ View constructive analysis as an **extension** of classical analysis.
- ▶ Formalize proofs in **TCF** (based on the Scott-Ershov model of partial continuous functionals), extract algorithms (in Minlog).
- ▶ **Data** are important (real number, continuous function ...).
- ▶ **Low type levels**: continuous $f: \mathbb{R} \rightarrow \mathbb{R}$ determined by its values on the rationals \mathbb{Q} .