

Uniform Spaces and Uniform Domains

--- Representing **a Uniform space + a Base**
for Computation ---

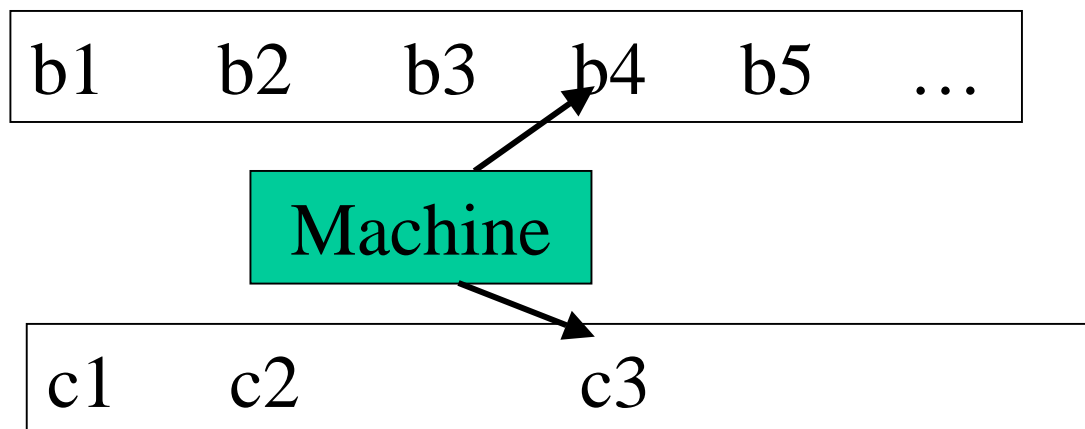
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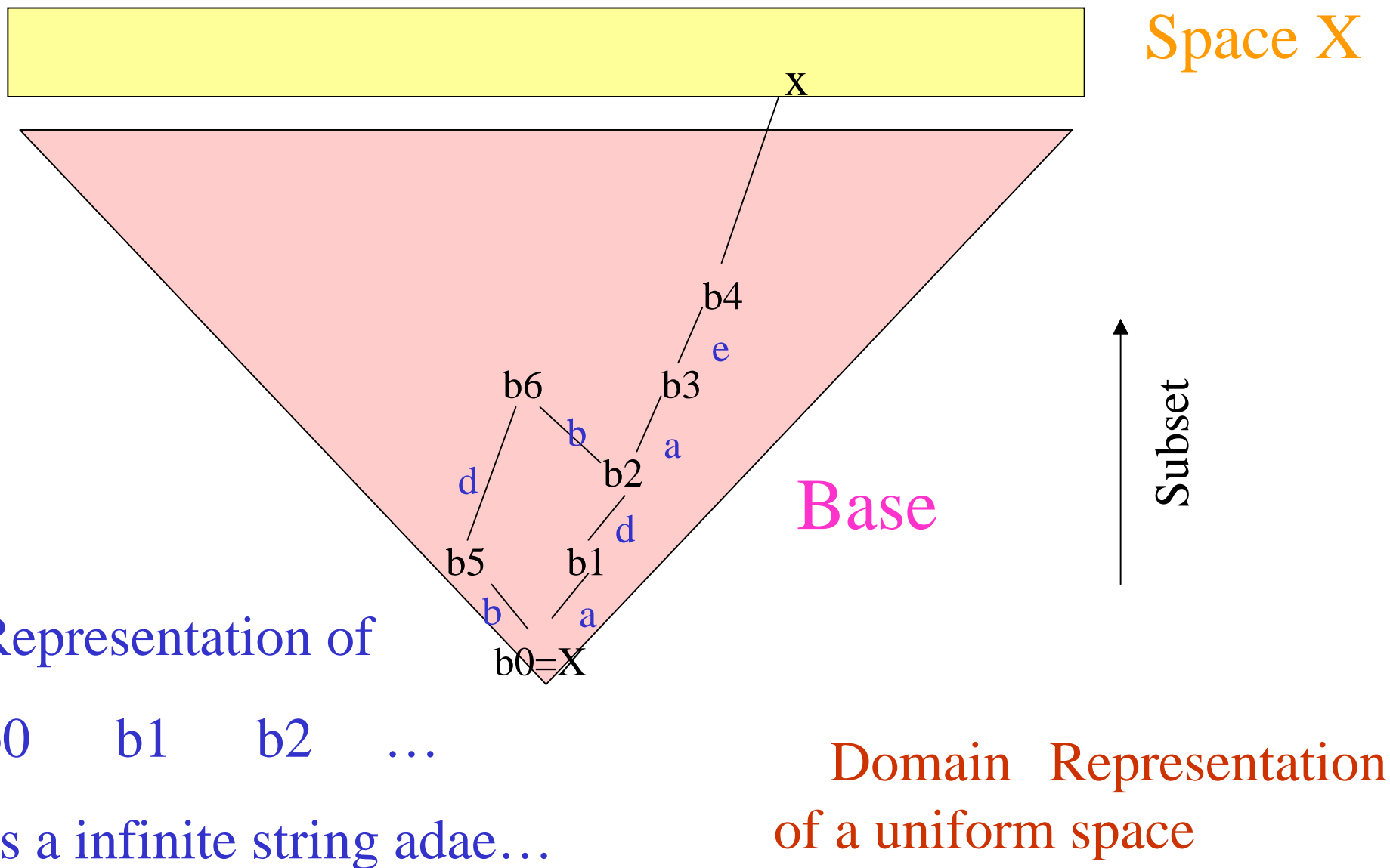
Computation over a Topological Space

[2nd-countable, T0 space]

- Fix a **base** and a representation of the base.
- Computation with respect to approximations.
- approximation ... given as a sequence of base elements
b1 b2 b3 b4 b5 ... x
- A machine which inputs/outputs infinite sequences of (representation of the) base.

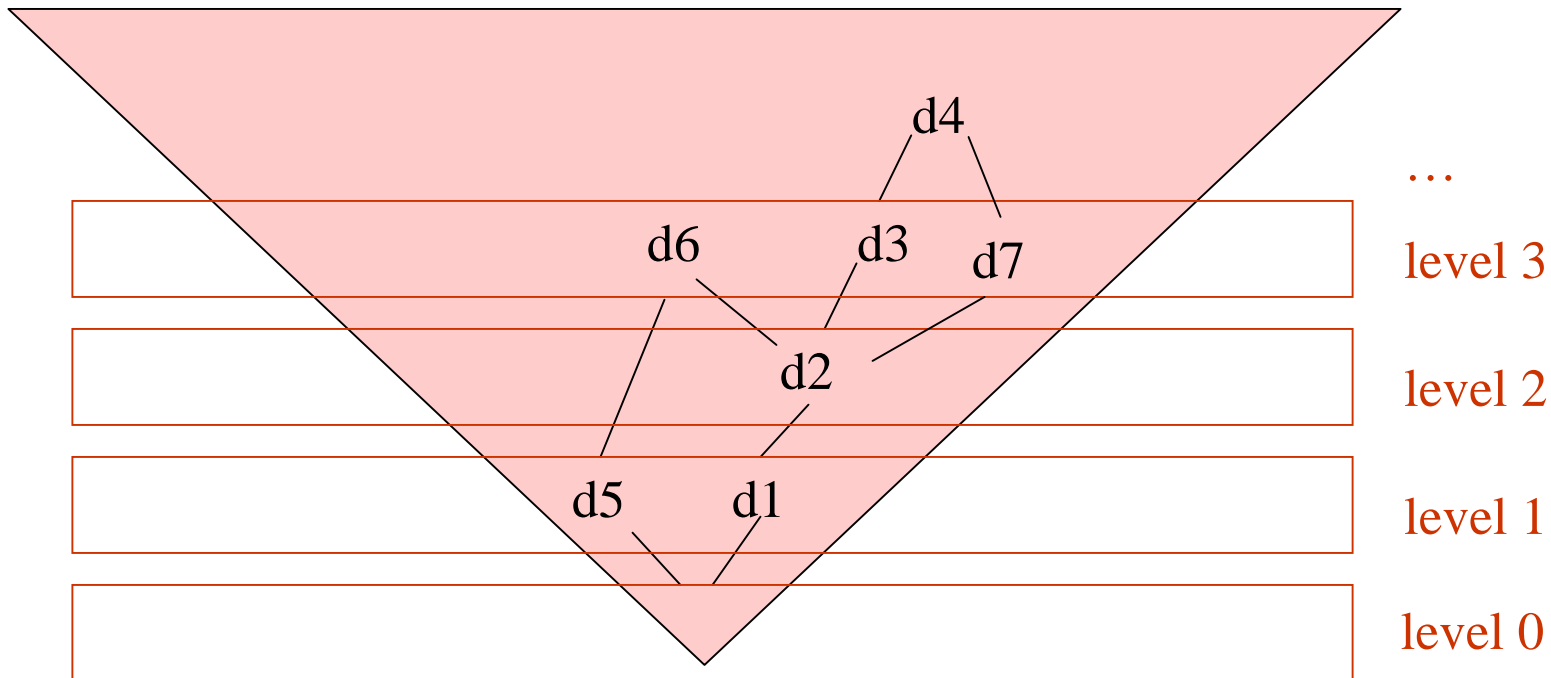


Base as a poset, points as limits



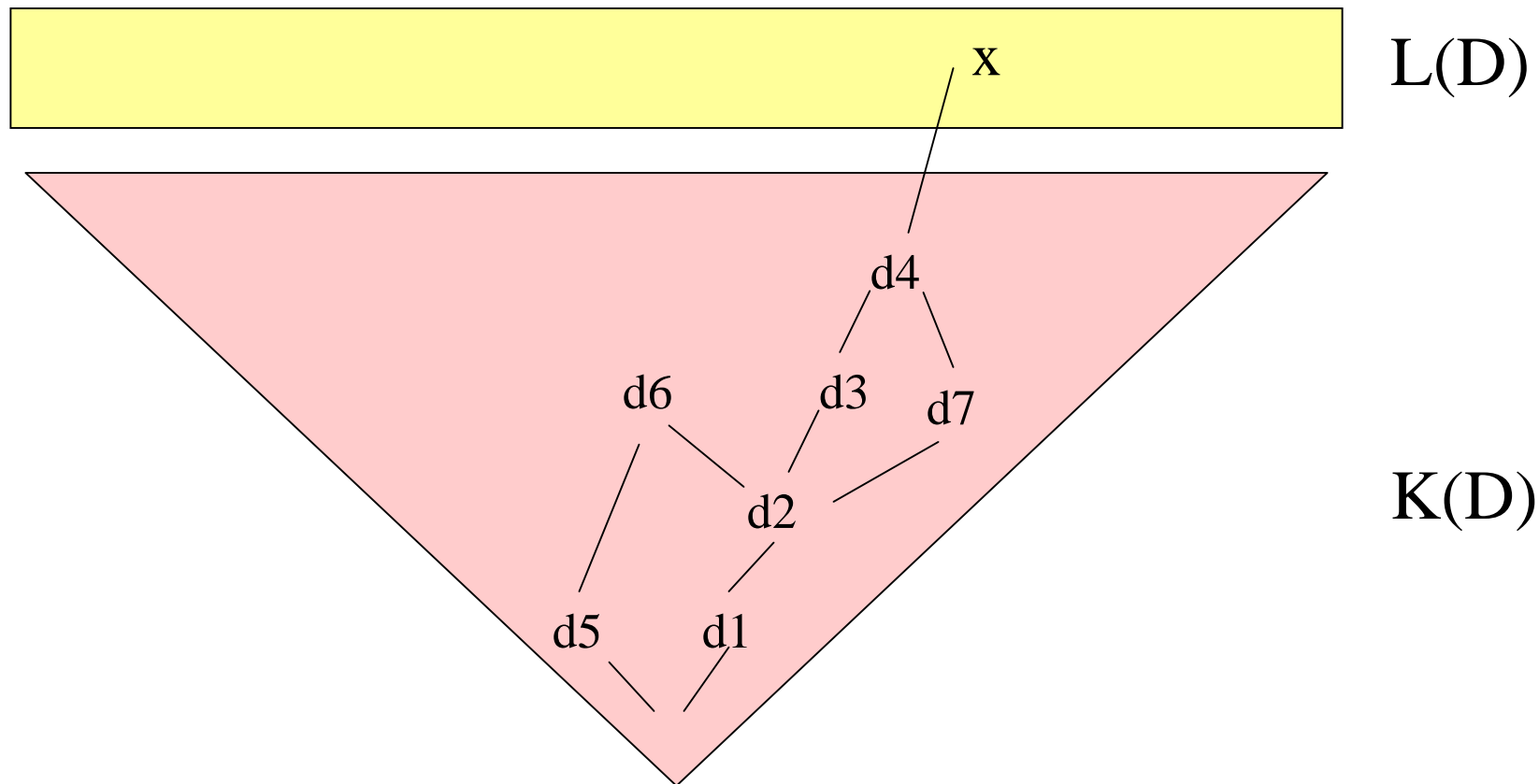
-type poset

- the level of d P : the maximal length of a chain $d_1 < d_2 < \dots < d$
- -type poset: every element $d \in P$ has a finite level.



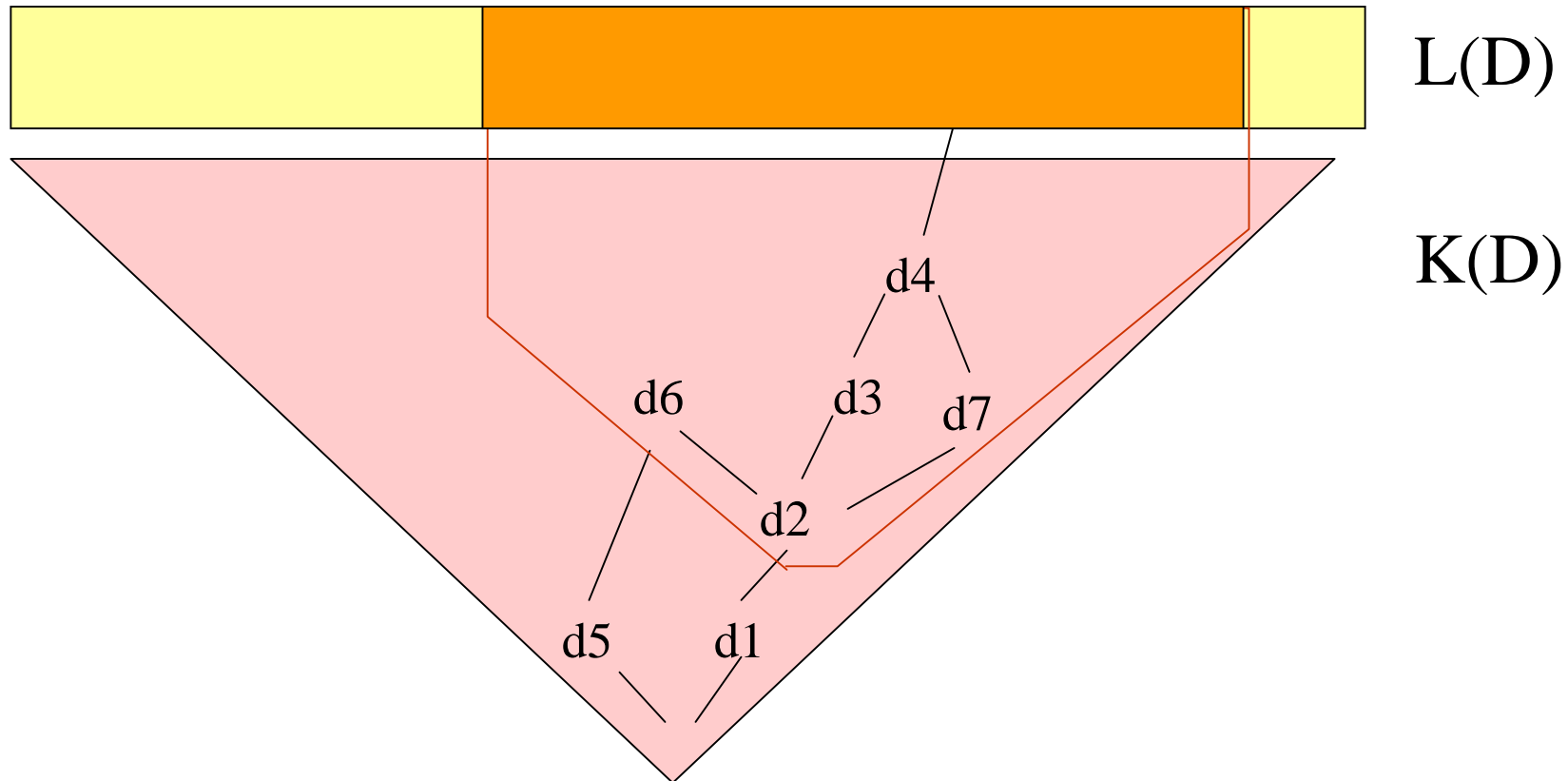
ω -type domain D

- Ideal completion of a ω -type poset
(the set of increasing sequences / equivalence)
- It is a ω -algebraic cpo.



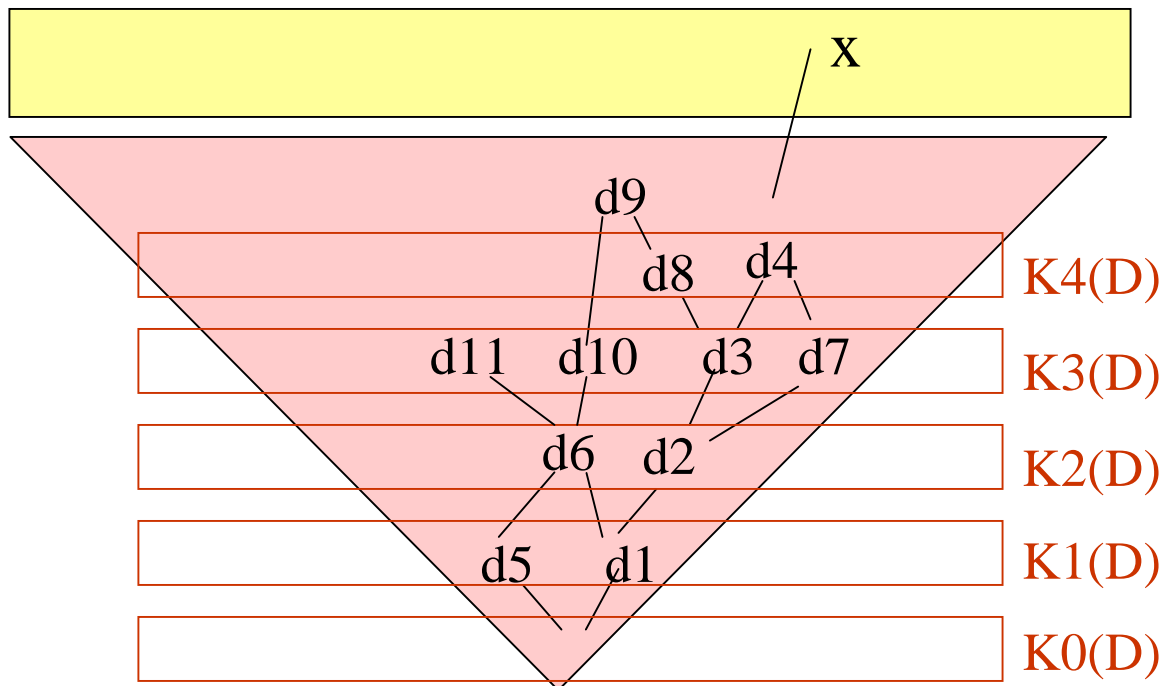
Scott Topology on a ω -type domain D

- (relative) Scott topology on $L(D)$
- $\{d \mid d \in K(D)\}$ as a base.
- T_0 space



Uniform Domain

- For $d \in \text{Kn}(D)$, define $d^* = \{e \in \text{Kn}(D) \mid e \leq d\}$
- Define $\text{mlb}(n)$ as the maximum level of a lower bound of d^* for $d \in \text{Kn}(D)$.
- Uniform domain is a Kn -type domain D such that $\text{mlb}(n) \leq n$ as $n \rightarrow \infty$.

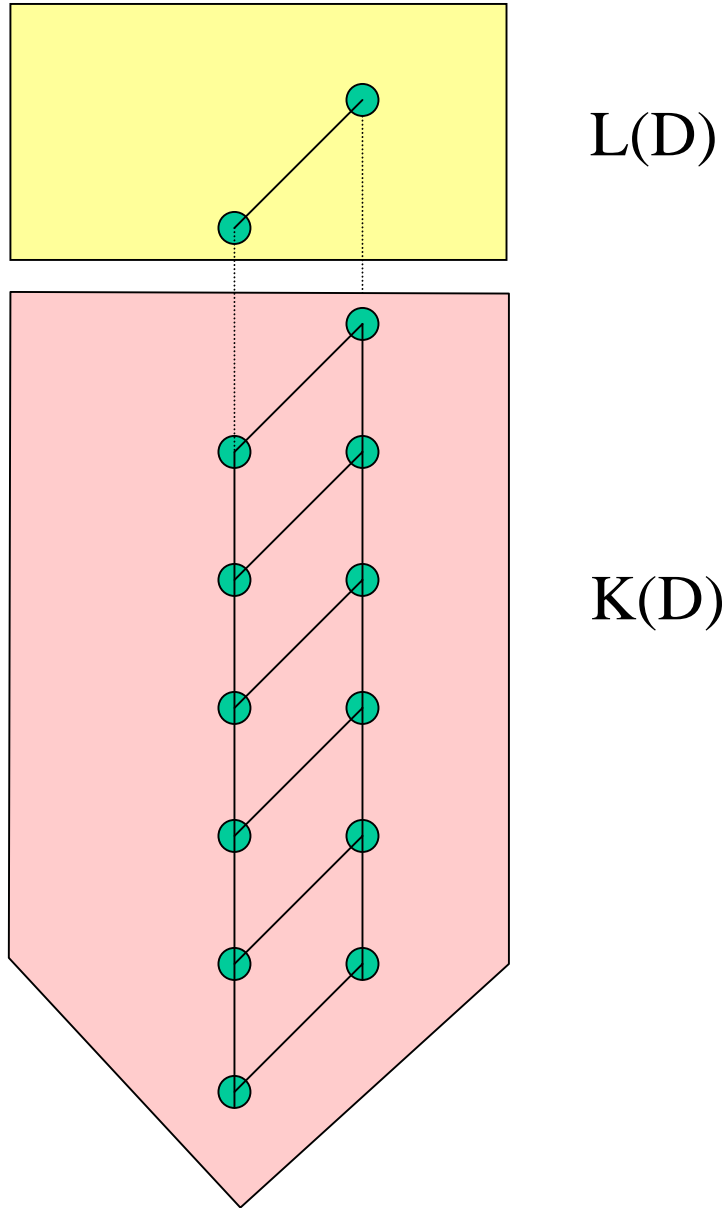


$L(D)$

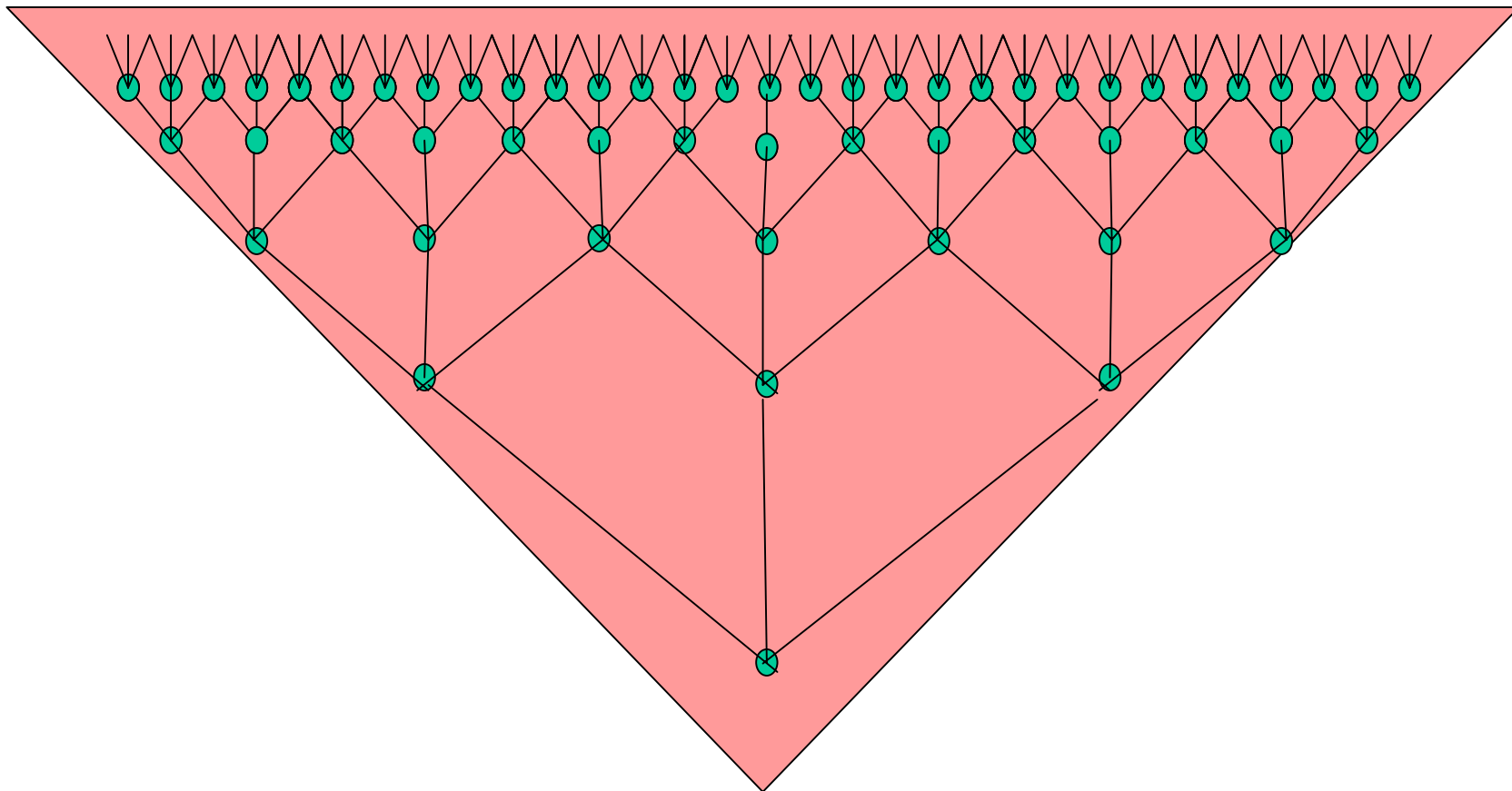
- $d3^* = \{d3, d7, d10\}$

- $\text{mlb}(3) = 1$

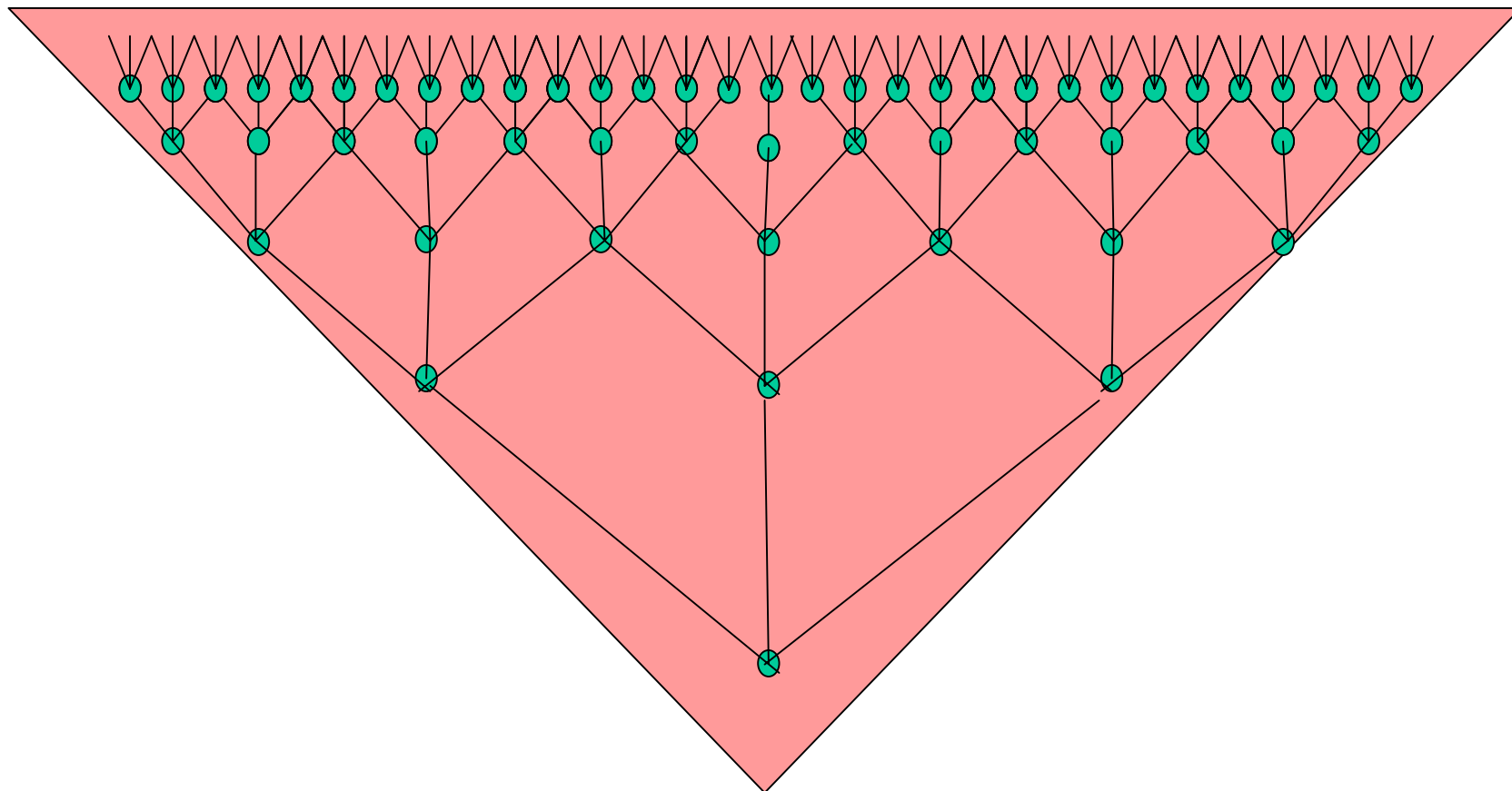
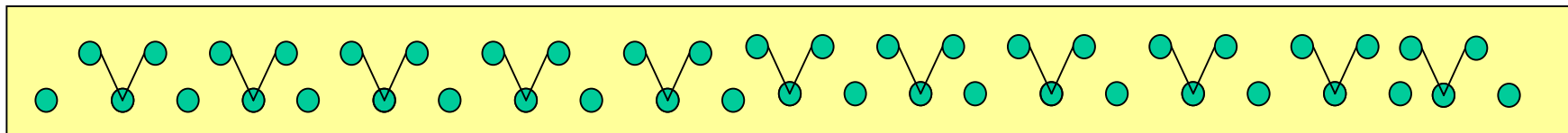
An example of a uniform domain



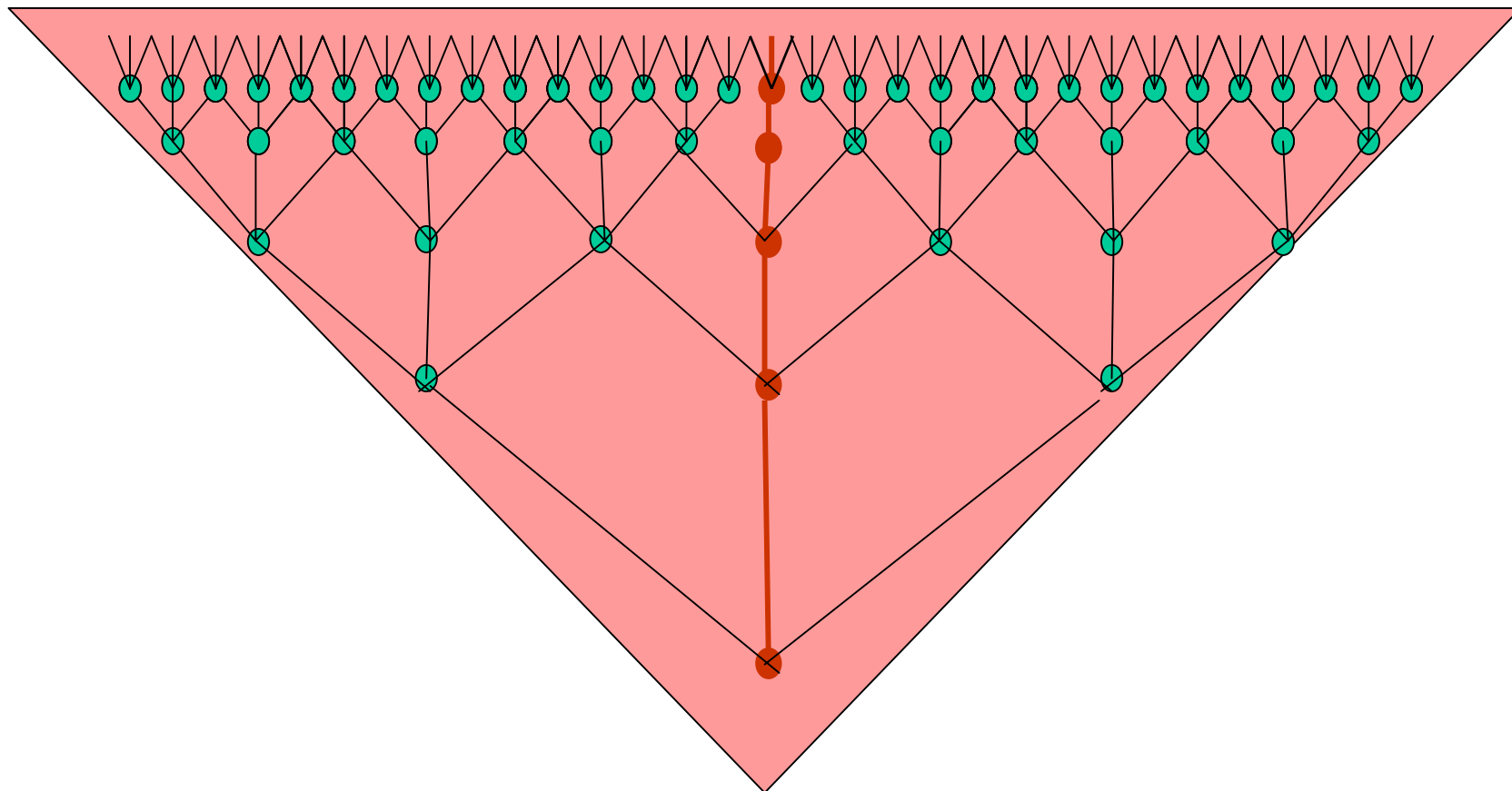
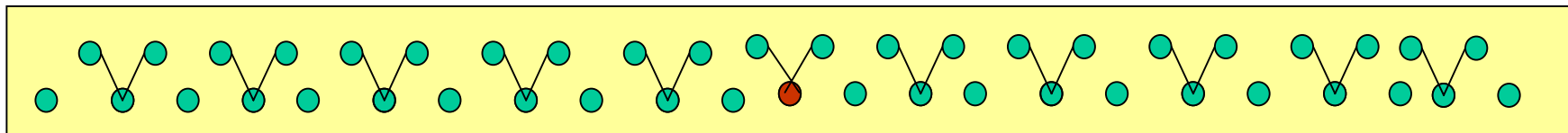
uniform domain RD



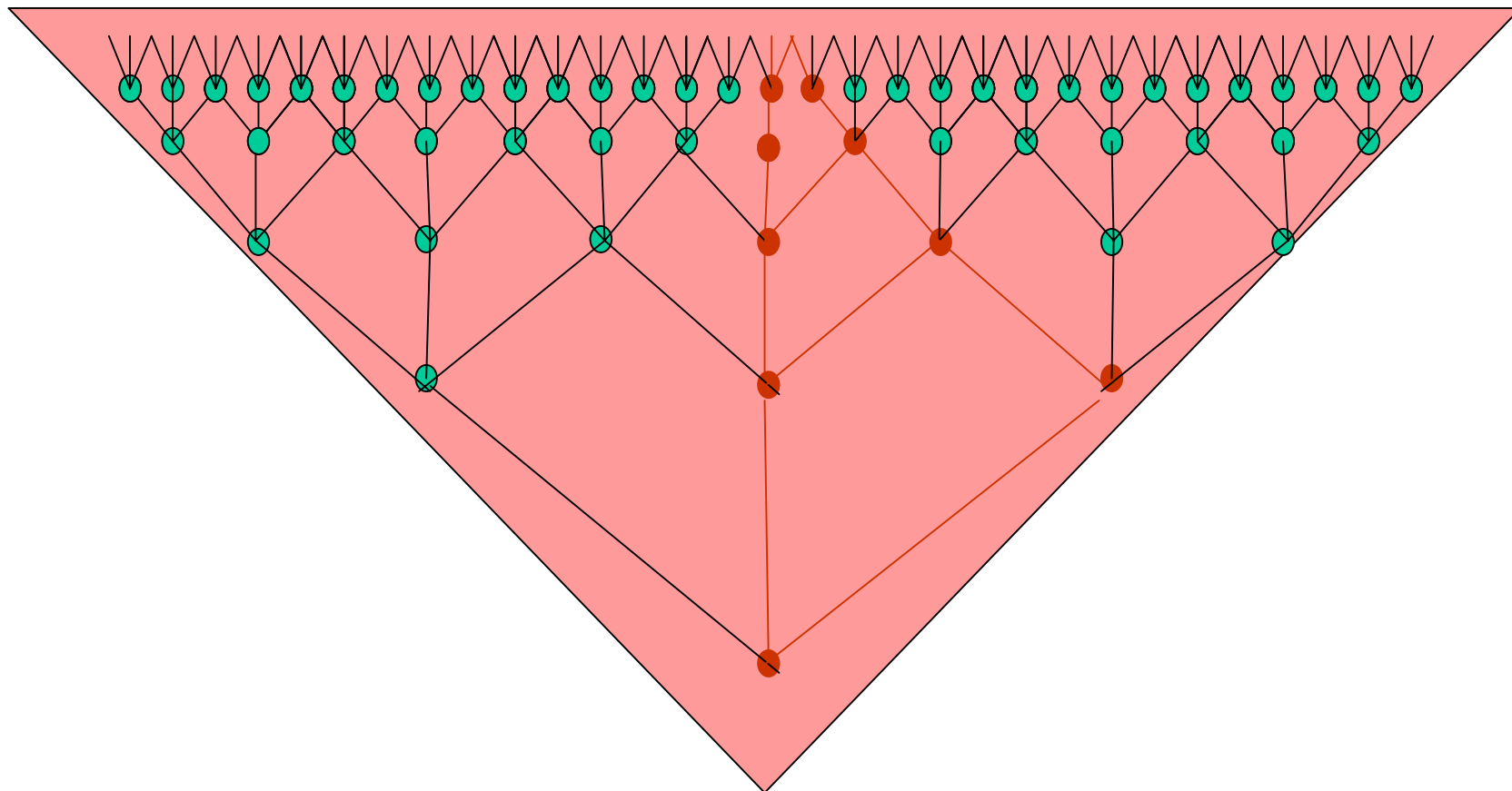
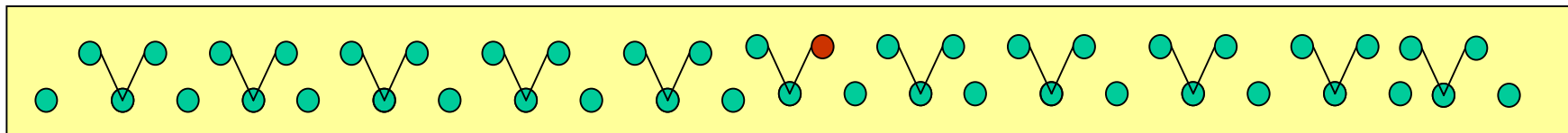
uniform domain RD



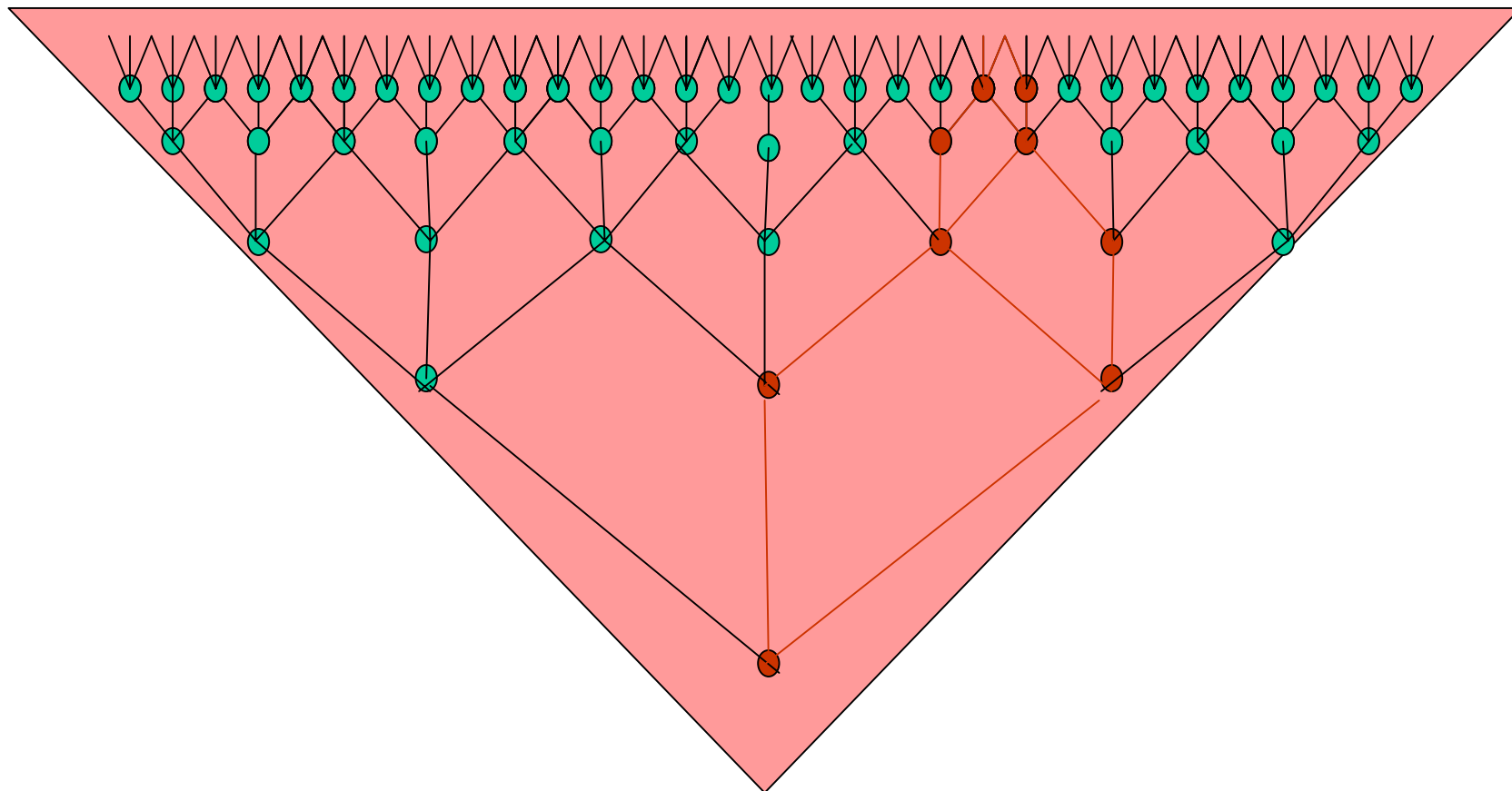
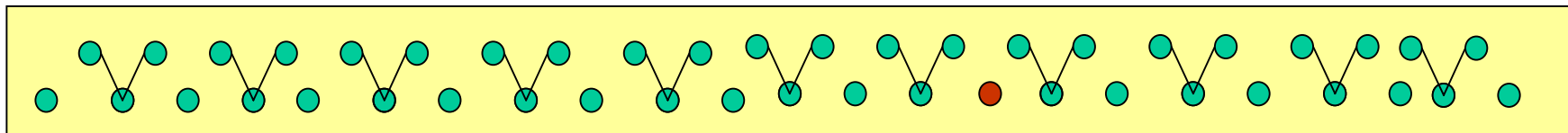
uniform domain RD



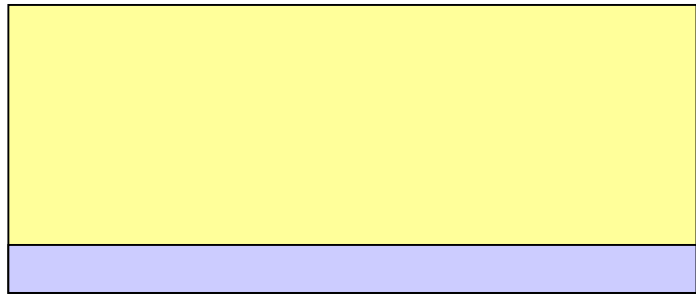
uniform domain RD



uniform domain RD

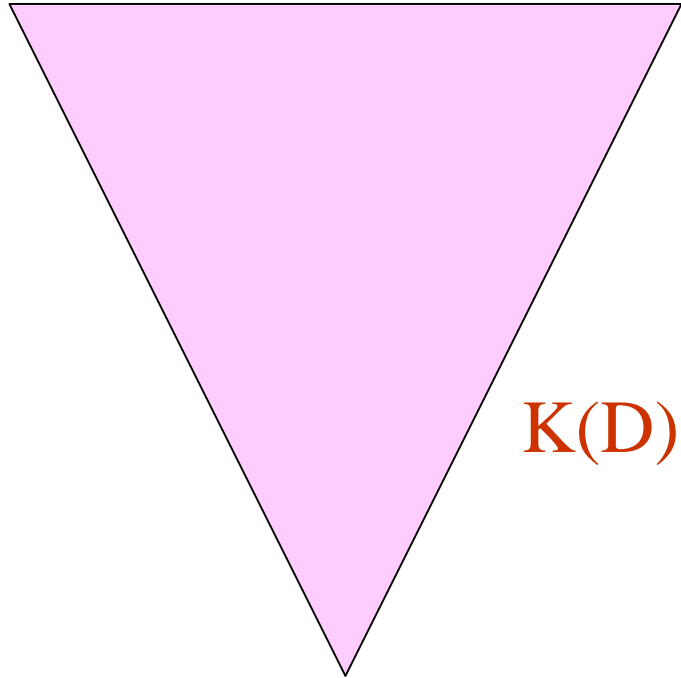


Minimal subspace of a uniform domain



$L(D)$

$M(D)$

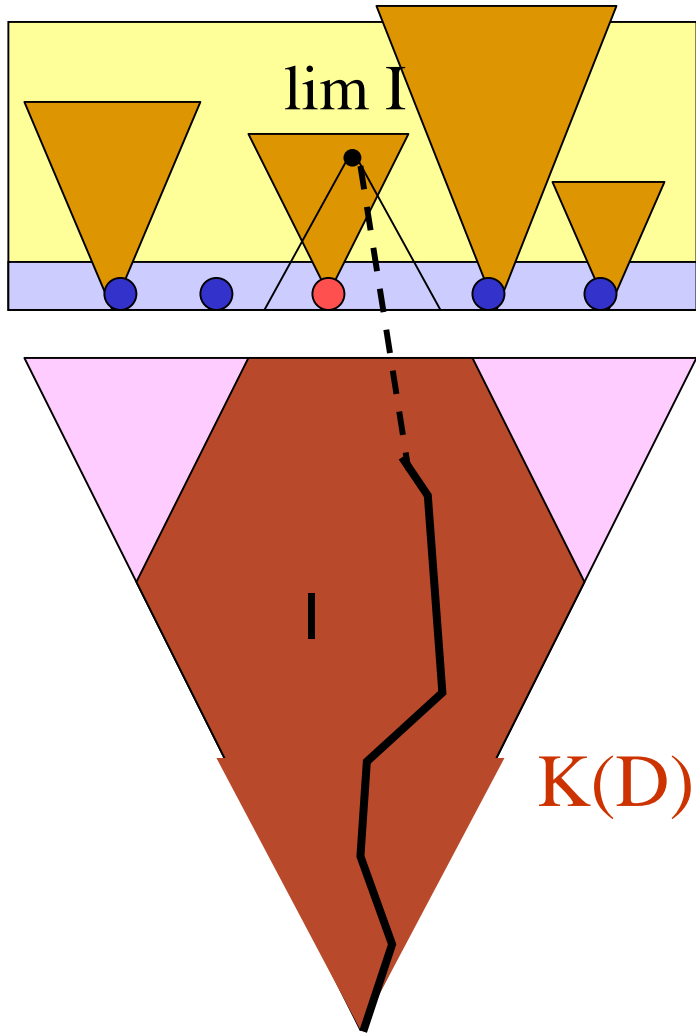


$K(D)$

Theorem When D is a uniform domain, $L(D)$ has a minimal subspace $M(D)$.

cf. $K(D)$ -type domain of $P(N)$
--- $K(D)$:finite set,
 $L(D)$:infinite set

Minimal subspace of a uniform domain (2)



$L(D)$

$M(D)$

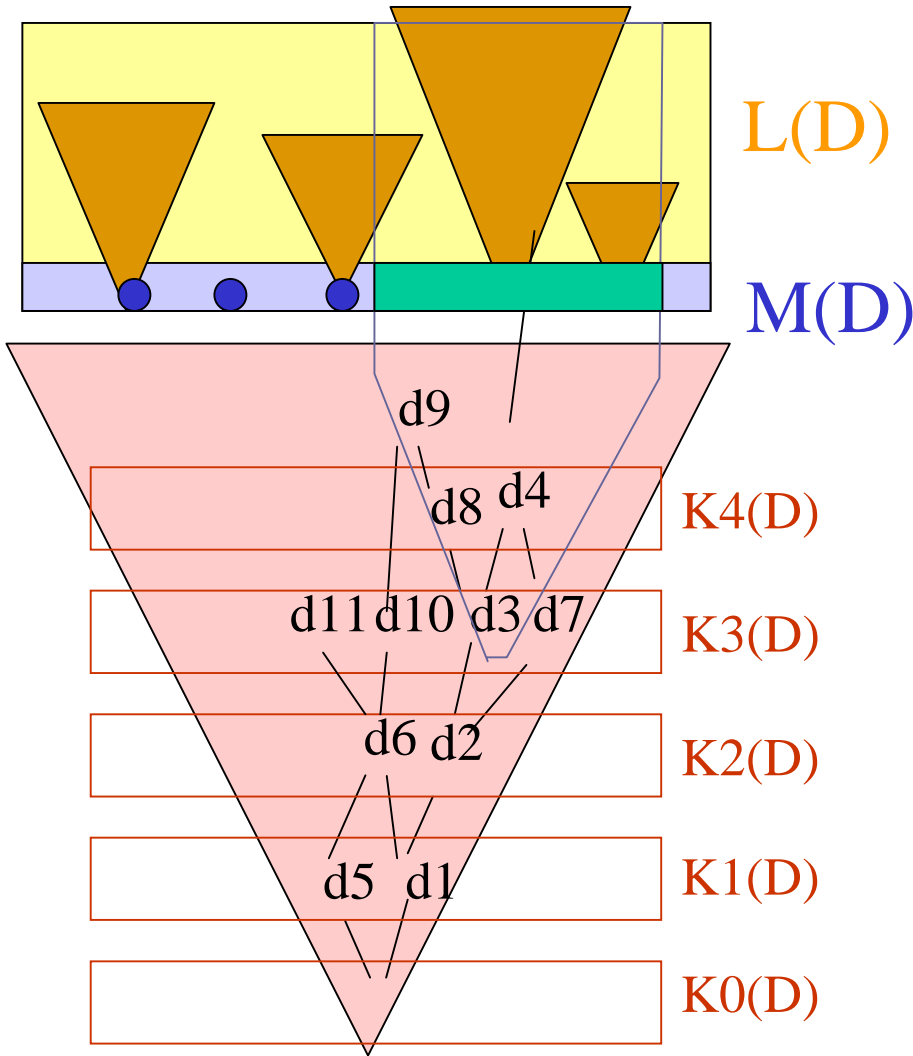
$K(D)$

Theorem When D is a uniform domain,

- $M(D)$ is a retract of $L(D)$.
- $M(D)$ is a Hausdorff space.

Each increasing sequence in $K(D)$ specifies one point of $M(D)$.

Uniform domain and uniform space.



- $[d] = d \in M(D)$
- $L(D)$ open set of $M(D)$
- $\mathcal{V}_n = \{ [d] \mid d \in K_n(D) \}$ open covering of $M(D)$

Theorem When D is a uniform domain, $M(D)$ is a countable complete uniform space with \mathcal{V}_n uniform coverings.

Cor. $M(D)$ is metrizable.

Countable complete uniform space with a base (X, μ)

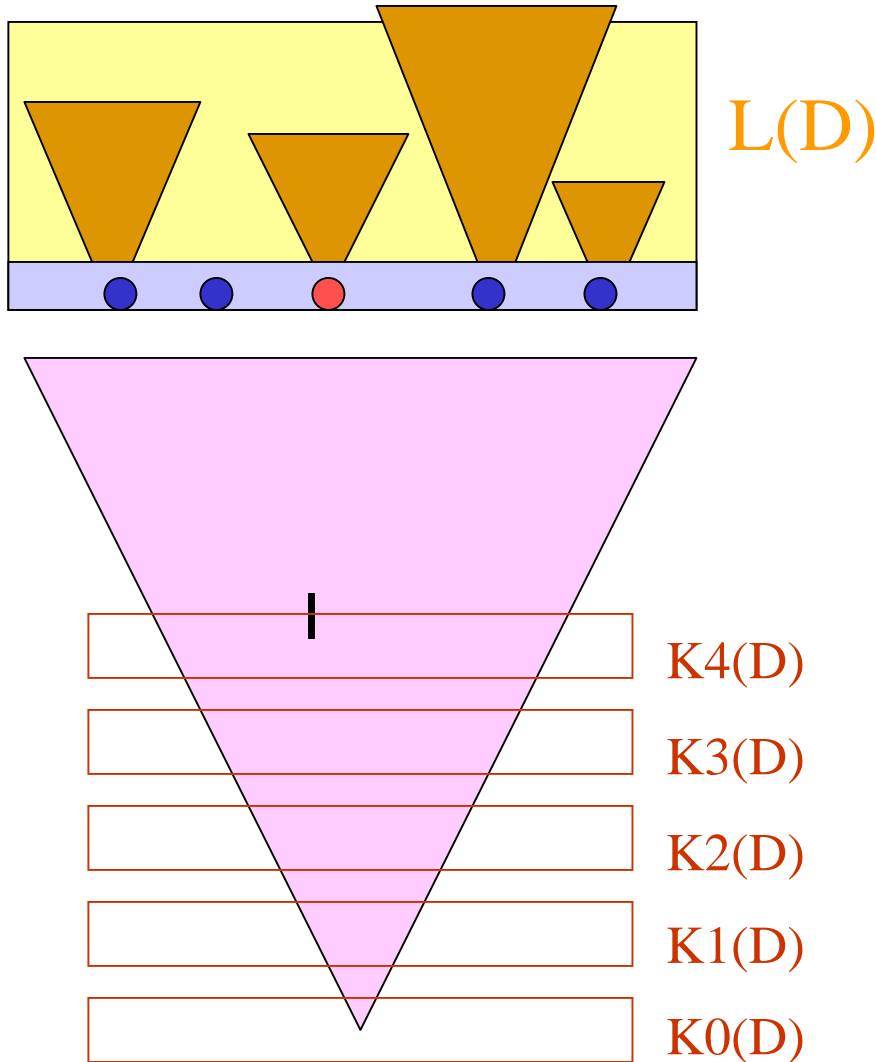
- X : topological space,
- $\mu = \{U_0, U_1, U_2, \dots\}$
A sequence of uniform open coverings s.t.
 U_n is a refinement of U_m for $n > m$,
 $U_n * U_m$ if $U_n \subseteq \{St(v, U_m) \mid v \in U_m\}$
- $St(A, U)$ is $\{u \in U \mid u \in A\}$

μ becomes a base of a uniformity on X

Theorem When (X, μ) is a countable complete uniform space with a base, we can form from μ a uniform domain D s.t. $M(D) = X$.

Totally bounded uniform domain

$K_n(D)$ finite set for all n



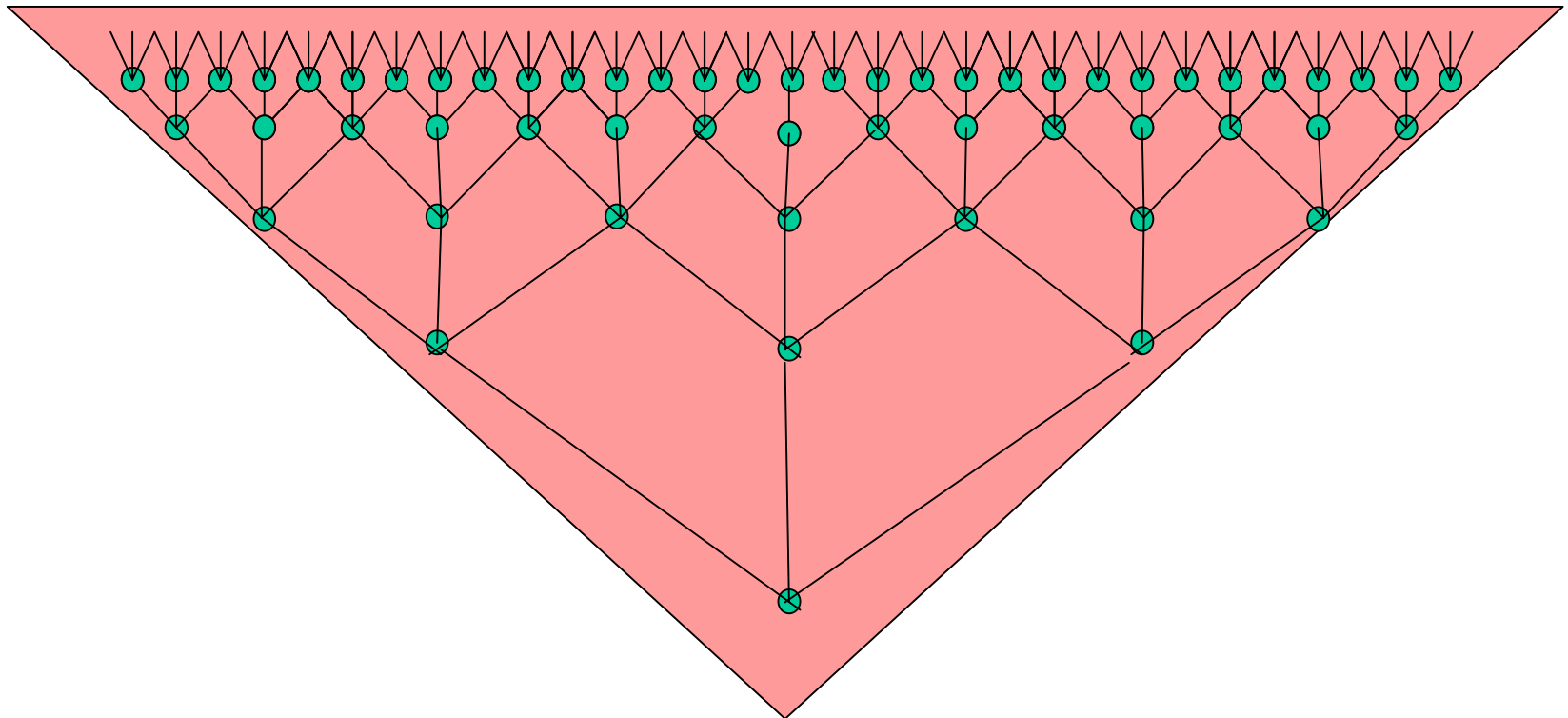
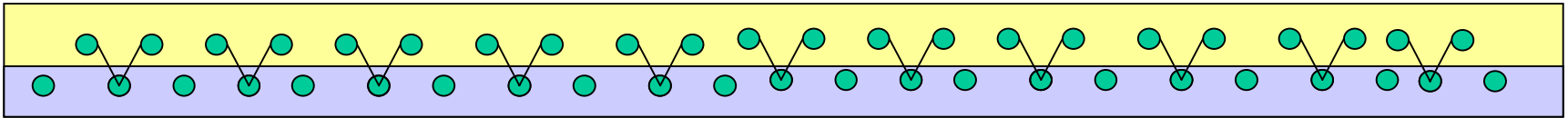
Theorem When D is a totally bounded uniform domain, $M(D)$ is a compact uniform space.

Theorem When (X, μ) is a compact countable uniform space, the corresponding uniform domain is totally bounded.

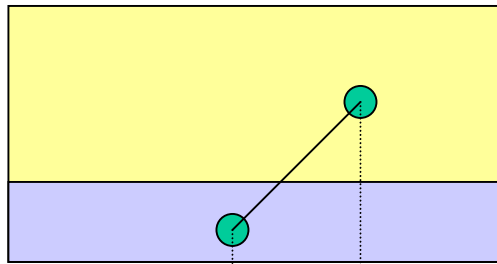
$M(\text{RD})$ is homeomorphic to $I=[0,1]$

Signed digit representation [Gianantonio]

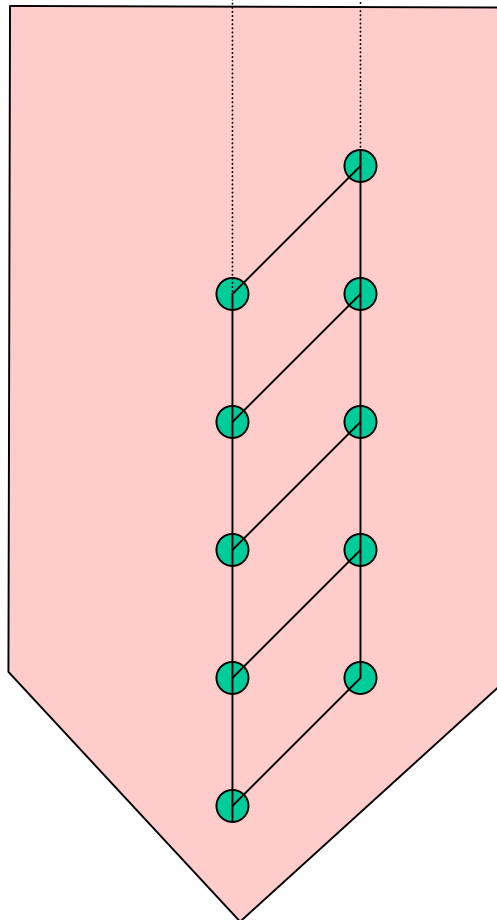
Gray code [Tsuiki]



An example $M(D)$ not dense in D



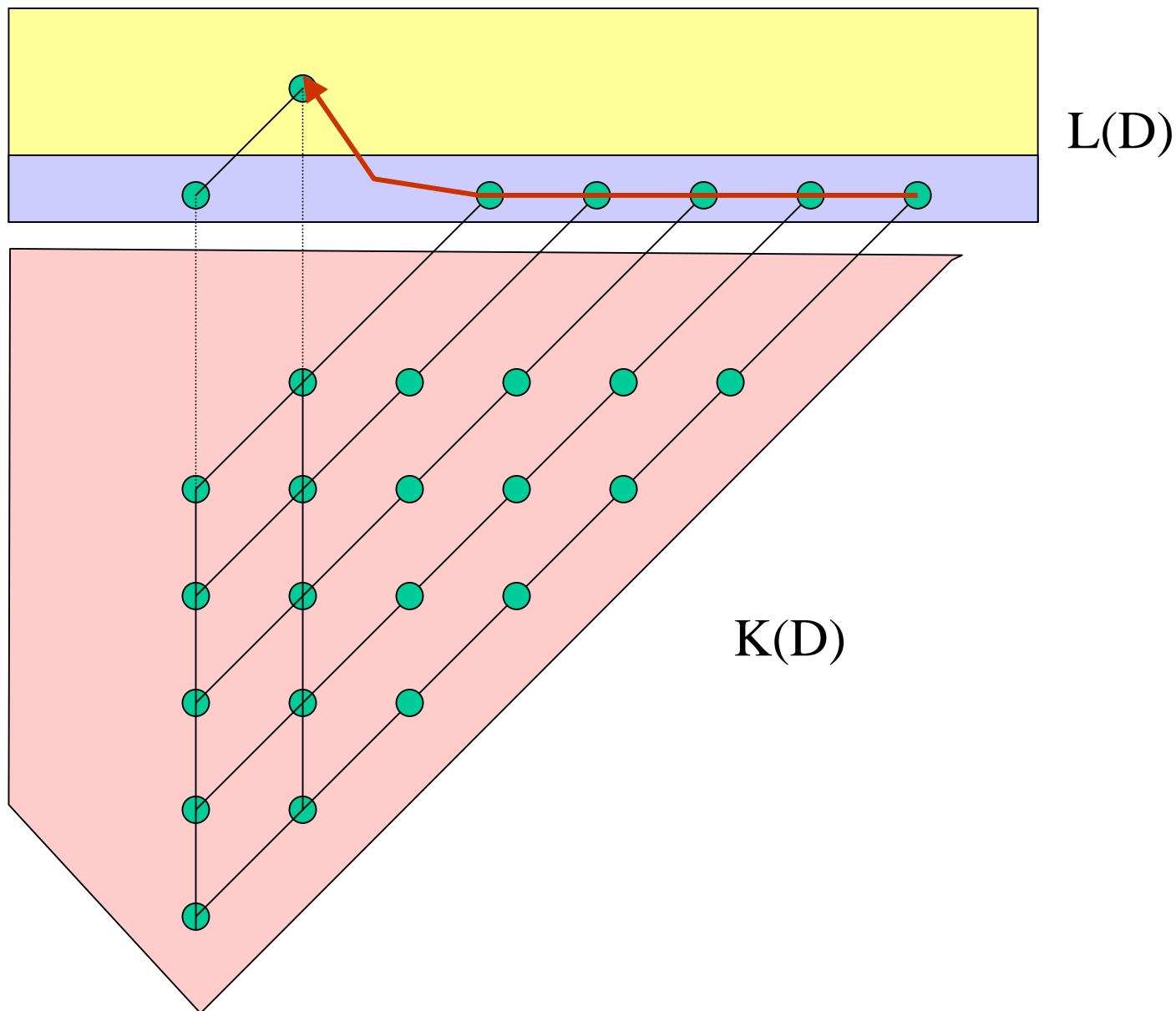
$L(D)$



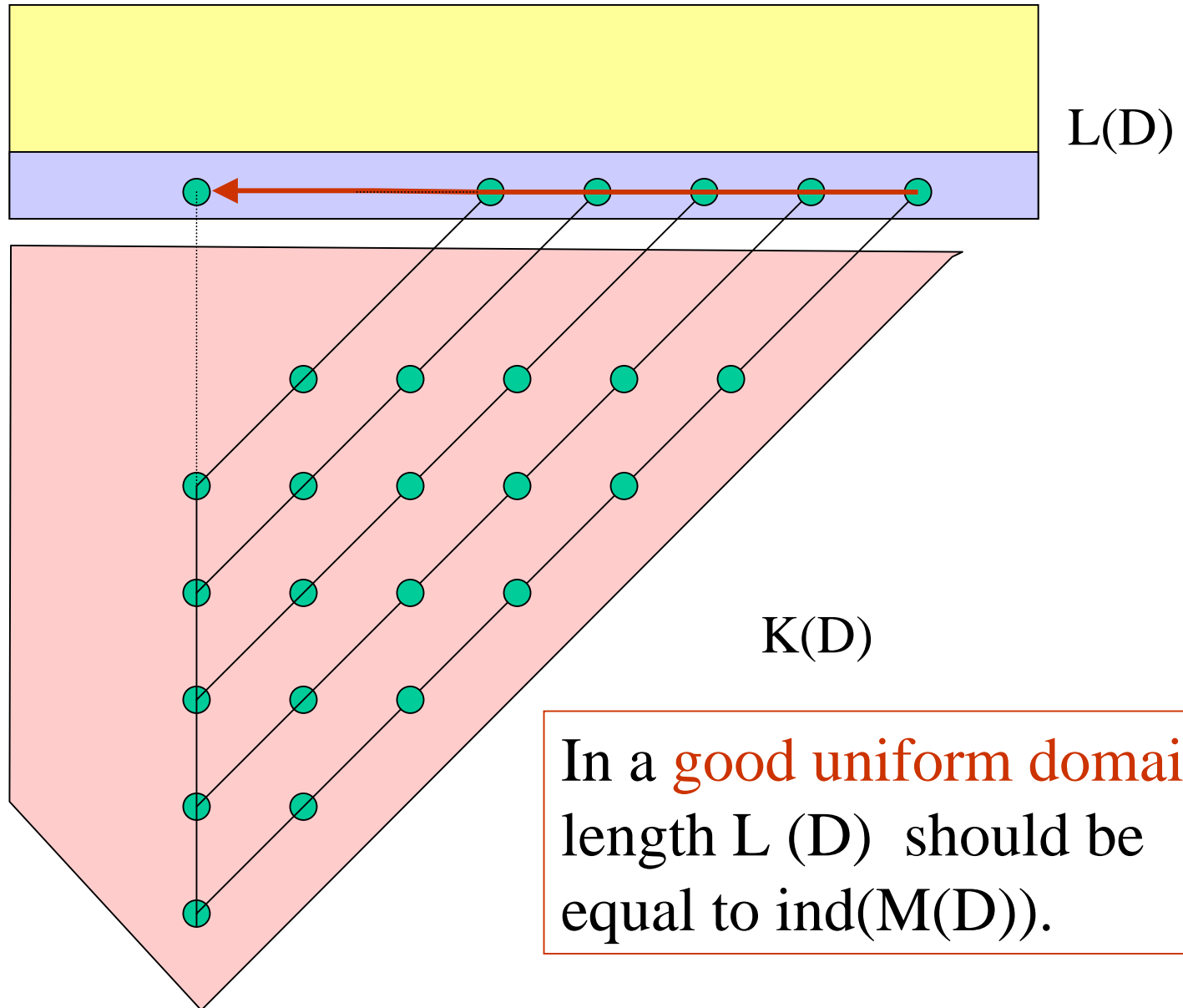
$K(D)$

In a **good uniform domain**,
 $M(D)$ should be dense in
 D

A uniform domain $M(D)$ dense in D



Another uniform domain with the same $M(D)$



In a **good uniform domain**, length $L(D)$ should be equal to $\text{ind}(M(D))$.

What is a good uniform domain?

- **Uniform domain = uniform space + base.**
- Which uniform domain is natural as a base of computation?
 - $M(D)$ dense in D . (d . e . d e^*)
 - $d \leq e$ iff $[d] \leq [e]$.
 - $\text{ind}(M(D)) = \text{length}(L(D))$

we have such domain for each separable metric space X [tsuiki,2002]

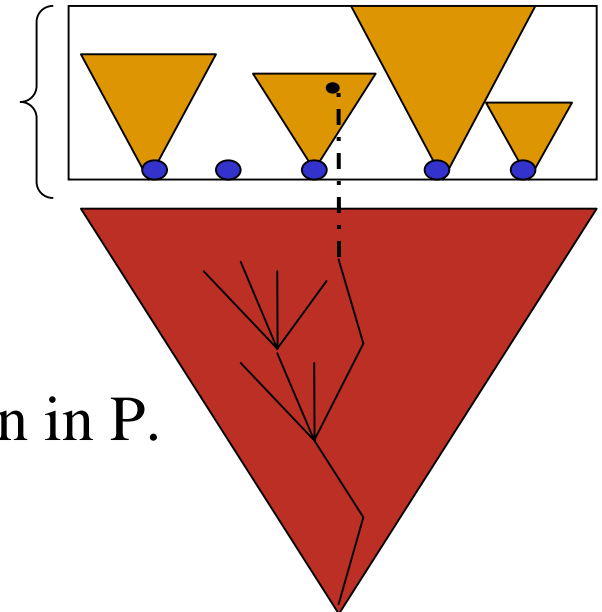
- In general we have

$$\text{ind}(L(D)) = \text{length}(L(D))$$

$$\text{ind}(M(D)) \leq \text{length}(L(D))$$

$\text{length}(P)$: the maximal length of a chain in P .

ind : Small inductive dimension.



Conclusion

- Computation ---- depends on the selection of a base

**Uniform domain =
Uniform space + base**

- Which property of a uniform domain is not dependent of the selection of a base?
- Are there any good conditions on a uniform domain which guarantees good computational properties?